Proof-theoretic aspects of obtaining a constructive version of the mean ergodic theorem

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Logic Colloquium '07, Wroclaw

'Proof mining' is the subfield of mathematical logic that is concerned with the extraction of additional information from proofs in mathematics and computer science.

G. Kreisel: What more do we know if we have proved a theorem by restricted means other than if we merely know the theorem is true?

 \Rightarrow Kreisel's unwinding program/proof mining.

Develop methods to unwind the computational content of ineffective proofs, i.e. proofs using full classical logic.

Quantitative information: Effective realizers and bounds (algorithms), complexity of realizers and bounds.

Qualitative information: Uniformities (bounds independent of certain parameters), weakening/elimination of premises.

Aims of proof mining: Classify theorems and proofs from which additional information can be extracted. Carry out case studies, i.e. analyse actual proofs in mathematics.

Methods of proof mining: Proof interpretations. Transform a proof P of a theorem A into an enriched proof P' of an equivalent theorem A' from which the desired information can be read off.

Applications of proof mining: E.g. algebra, analysis (fixed point theory and approximation theory), combinatorics, number theory and computer science - and recently: ergodic theory.

Why proof mining in ergodic theory?

- Ergodic theory often uses abstract, non-computational structures and techniques.
- Ergodic theory is often concerned with asymptotic behaviour of iterative processes \Rightarrow extraction of bounds.
- Ergodic theory has connections with many other areas of mathematics, e.g. combinatorics and number theory.

Overview

Introduction

- Mean Ergodic Theorem
- Metatheorems
- Proof Analysis

Ergodic Theory

Very informal introduction to ergodic theory:

- **•** Take a measure space (X, \mathcal{B}, μ) .
- Consider a measure preserving map $T: X \to X$, i.e. $\mu(TA) = \mu(A)$ for all $A \in \mathcal{B}$.
- Study the asymptotic behaviour of T^nA .

One can study a measure preserving system by studying the Hilbert space $L_2(X, \mathcal{B}, \mu)$ and the induced operator U_T .

Mean Ergodic Theorem

Let $(X, \langle \cdot, \cdot \rangle)$ be a Hilbert space, $T : X \to X$ a nonexpansive mapping and for $f \in X$ define $A_n f := \frac{1}{n+1} \sum_{i=0}^n T^i f$.

Mean Ergodic Theorem: The sequence $A_n f$ converges in the Hilbert space norm.

Questions: Does a computable rate of convergence exist? If not, what kind of rates can we obtain? In what way do those depend on f, T and the space $(X, \langle \cdot, \cdot \rangle)$?

Mean Ergodic Theorem

No full rate of convergence, even for space $L_2(X, \mathcal{B}, \mu)$ counterexamples using the halting problem. Full rate of convergence for $L_2(X, \mathcal{B}, \mu)$ if *T* is ergodic.

Consider classically equivalent, no-counterexample version of the mean ergodic theorem:

Mean Ergodic Theorem (n.c.i.): For every $f \in X$, nonexpansive $T: X \to X$, $M: \mathbb{N} \to \mathbb{N}$ and $\varepsilon > 0$, there is an $n \in \mathbb{N}$ such that $||A_m f - A_n f|| \le \varepsilon$ for all $m \in [n, M(n)]$.

From a standard proof of the mean ergodic theorem, we extract effective bounds for this no-counterexample version.

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Metatheorems

In what kind of formal system can we prove the mean ergodic theorem? What kind of additional information do the metatheorems predict?

Definitions:

- Finite types T: (i) $\mathbb{N} \in \mathbf{T}$, (ii) $\rho, \sigma \in \mathbf{T} \Rightarrow \rho \rightarrow \sigma \in \mathbf{T}$.
- PA $^{\omega}$: Peano Arithmetic in all finite types.
- DC: Axiom schema of dependent choice.
- $\mathcal{A}^{\omega} := \mathsf{P}\mathsf{A}^{\omega} + \mathsf{D}\mathsf{C}$ classical analysis in all finite types.

Metatheorems

Using a monotone variant of Gödel's functional ('Dialectica') interpretation, one may prove general metatheorems about the extraction of effective, uniform bounds from proofs of $\forall \exists A_{qf}$ statements in \mathcal{A}^{ω} (U.Kohlenbach).

Large parts of classical analysis can be formalized in \mathcal{A}^{ω} , in particular concerning complete seperable metric spaces.

What about e.g. abstract Hilbert spaces $(X, \langle \cdot, \cdot \rangle)$?

Define the following formal system $\mathcal{A}^{\omega}[X, \langle \cdot, \cdot \rangle]$:

- **J** Let \mathbf{T}^X be the finite types over \mathbb{N} and new type X.
- Extend \mathcal{A}^{ω} to finite types \mathbf{T}^X .
- New constants: $0_X, 1_X, +_X, \cdot_X$ and $\langle \cdot, \cdot \rangle_X$.
- New axioms: defining axioms for Hilbert spaces.

G./Kohlenbach(TAMS,2005/to appear): General logical metatheorems about the extraction of effective uniform bounds from proofs in $\mathcal{A}^{\omega}[X, \langle \cdot, \cdot \rangle]$ (and similar theories).

Definition: We call the following finite types over \mathbb{N} , X *small*: \mathbb{N} , X, $\mathbb{N} \to \mathbb{N}$, $\mathbb{N} \to X$, $X \to X$.

Definition: A formula F_{\forall} , i.e. F_{\exists} is a formula $\forall \underline{x}^{\underline{\sigma}} F_{qf}(\underline{x})$, i.e. $\exists \underline{x}^{\underline{\sigma}} F_{qf}(\underline{x})$, with F_{qf} quantifier-free and all types σ_i small.

Definition: Let $(X, \langle \cdot, \cdot \rangle)$ be a Hilbert space. An operator $T: X \to X$ is nonexpansive – short T n.e. – if

$$\forall f, g \in X(\|Tf - Tg\| \le \|f - g\|).$$

Corollary^a: Let $(X, \langle \cdot, \cdot \rangle)$ be a Hilbert space. If

$$\forall f^X, T^{X \to X}, k^0, M^1(T \text{ n.e.} \land \\ \forall u^0 B_{\forall}(f, T, k, M, u) \to \exists v^0 C_{\exists}(f, T, k, M, v)),$$

is provable in $\mathcal{A}^{\omega}[X, \langle \cdot, \cdot \rangle]$, then there is a computable $\varphi : \mathbb{N} \times \mathbb{N} \times \mathbb{N}^{\mathbb{N}} \to \mathbb{N}$, so that

$$\forall f^X, T^{X \to X}, k^0, M^1(||f||, ||f - Tf|| \le b \land T \text{ n.e. } \land \\ \forall u^0 \le \varphi(b, k, M) B_\forall \to \exists v^0 \le \varphi(b, k, M) C_\exists)$$

holds in every Hilbert space $(X, \langle \cdot, \cdot \rangle)$.

^{*a*}I.e., corollary to most general metatheorems in TAMS-paper.

Easy to check that the mean ergodic theorem can be proved in $\mathcal{A}^{\omega}[X, \langle \cdot, \cdot \rangle]$ and (the no-counterexample version) has the suitable, logical form for the metatheorem.

Metatheorem predicts (i.e. guarantees) effective bounds depending only on ||f||, M and ε . Bounds are independent of the space $(X, \langle \cdot, \cdot \rangle)$ and the mapping T, and uniform on every norm-bounded, *not necessarily compact* ball in X.

Joint work with J.Avigad and H.Towsner: Extract effective bounds from standard proof of mean ergodic theorem.

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Proof Analysis

Standard proof of mean ergodic theorem:

$$U = \overline{\{u - Tu | u \in X\}}, \ V = \{v \in X | v = Tv\}, \ U \perp V.$$

$$||A_n(u - Tu)|| \le \frac{2||u||}{n}$$
, $A_nv = v$, $X = U \oplus V$.

Let $f \in X$ be given, let $f_U = u - Tu$ be close to the projection of f on U and let $f^* = f - f_U$.

Since f_U is close, $||A_n f^*||$ remains stable for all $n \in \mathbb{N}$. Let N be large enough then $||A_n f_U||$ is small for all $n \ge N$.

Then $||A_n f - A_m f||$ small for all $n, m \ge N$.

Proof Analysis

Making this argument effective requires:

- Solution Rate of convergence for some sequence $g_n = u_n Tu_n$ of approximations to projection of f onto U.
- Bounds on the norms of elements u_n .

We will carry out the effective proof as if we had a rate of convergence for g_n , but will be able to obtain bounds using a slightly weaker notion of convergence for g_n .

Recall: We only want to find an interval [n, M(n)] where $A_n f$ is stable (within a given ε).

Proof Analysis - bounds on $||u_n||$

Observation: Sufficient to consider projection of f onto $U_f = \overline{\{T^i f - T^{i+1} f | i = 0, 1, ...\}}$, as $A_n f$ lies in the span of $\{f, Tf, T^2 f, ...\}$.

Define

$$u_0 = \frac{\langle f, f - Tf \rangle}{\|f - Tf\|^2} f, \quad u_{n+1} = u_n + \frac{\langle f - g_n, T^n f - T^{n+1} f \rangle}{\|T^n f - T^{n+1} f\|^2} T^n f,$$

then $g_n = u_n - Tu_n$ approximates projection of f onto U_f .

To bound $||u_i||$ we need lower bound on $||T^if - T^{i+1}f||$.

Proof Analysis - bounds on $||u_n||$

Second observation: If ||f - Tf|| is small, $||A_nf||$ is stable for a long time (using triangle inequality).

Either ||f - Tf|| is small enough to get $||A_{M(0)}f - f|| \le \varepsilon$. Or we have a lower bound on ||f - Tf|| and thus an upper bound on $||u_0||$.

Similar, for general $||u_n||$, i.e. we have a sequence of upper bounds for $||u_n||$ in terms of ||f||, ε and M.

Proof Analysis - convergence of g_n

Let u_n and g_n be as before. Assume we have a modulus of convergence for g_n . For any *i*

$$||A_{M(n)}f - A_nf|| \le ||A_{M(n)}(f - g_i) - A_n(f - g_i)|| + ||A_{M(n)}g_i|| + ||A_ng_i||.$$

By direct calculation we find a $\delta > 0$ s.t. if $||g_i - g_j|| \le \delta$ for all j > i, then $||A_{M(n)}(f - g_i) - A_n(f - g_i)|| \le \varepsilon/2$ for all n.

Using a modulus of convergence for g_n , we find such an *i*.

Making *n* large enough, we get $||A_{M(n)}g_i|| + ||A_ng_i|| \le \varepsilon/2$.

Proof Analysis - convergence of g_n

For all *n*, we have $||g_n|| \le ||g_{n+1}|| \le ||f||$. By the principle of convergence for bounded, monotone sequences, the sequence $||g_n||$ converges (and thus also g_n converges).

Interpreting (countable) choice to obtain a modulus of convergence for g_i requires bar-recursion. This would cause very complicated bounds for our constructive mean ergodic theorem.

Inspired by a technique (due to Kohlenbach) to eliminate certain simple instances of choice such as PCM, we observe the following:

Proof Analysis - convergence of g_n

We wrote: "find an *i* s.t. $||g_j - g_i|| \le \delta$ for all j > i".

In fact: We only need $||g_j - g_i|| \le \delta$ for a specific *j* given in terms of *i* and other parameters of the theorem. This yields a sequence of disjoint nonempty intervals $[i_k, j_k]$.

Using monotonicity and boundedness of $||g_n||$, we see that $||g_{j_k} - g_{i_k}|| \ge \delta$ only for finitely many $[i_k, j_k]$. Thus for one of those intervals we have the result.

Final quirk: The j_k depend on n, M(n) - but the number of intervals $[i_k, j_k]$ we need to consider is independent of n.

Proof Analysis - putting it all together

Solution: Define $i_0 = 0$, n_k in terms of i_k and i_{k+1} in terms of n_k ($j_k = i_{k+1}$):

$$i_0 := 0 \qquad n_k := \left\lceil \frac{b^2}{\varepsilon^2} \sum_{l=0}^{i_k} M(\frac{2lb}{\varepsilon}) \right\rceil,$$
$$i_{k+1} := i_k + \left\lceil \frac{2^{15}M(n_k)^4 b^4}{\varepsilon^4} \right\rceil.$$

Let $d = \frac{512b^2}{\varepsilon^2}$ and $N(b, k, n) = \frac{2n_d b}{\varepsilon}$, then

$$\forall f^X, T^{X \to X}, \varepsilon > 0, M : \mathbb{N} \to \mathbb{N} \exists n \le N (\|f\| \le b \land T \text{ n.e. } \land M(n) > n \to \|A_{M(n)}f - A_nf\| \le \varepsilon).$$

Conclusions

We have obtained a constructive proof of the following:

For every $f \in X$, $T: X \to X$ nonexpansive, $M: \mathbb{N} \to \mathbb{N}$, $\varepsilon > 0$ there exists an $n \le N$ – computable in ||f||, M and ε – such that $||A_m f - A_n f|| \le \varepsilon$ for every $m \in [n, M(n)]$.

- Proof theoretic guidelines to put proof (definitions, lemmas, theorem) into suitable logical form.
- Interpretation of lemmas: mostly direct computation.
- Elimination of instance of PCM.
- Extraction of computable bound as predicted by metatheorems.

References

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