

# ON MODELS OF PARACONSISTENT ANALOGUE OF THE SCOTT LOGIC

Marina Stukacheva

Sobolev Institute of Mathematics  
Novosibirsk, Russia

We consider the class JHN of all non-trivial extensions of the minimal (or Johans-son's) logic (denoted  $L_j$ ).

It is well-known that JHN partition into three disjoint subclasses [Odintsov 1998].

There are

the class INT of intermediate logics,

the class NEG of negative logics containing formula  $\perp$ , and

the class PAR of proper paraconsistent extensions of  $L_j$  consisting of logics not belonging to the first two classes.

In 1989 Zakharyashchev associated with every finite rooted intuitionistic frame  $\mu = \langle W, R \rangle$  a formula such that it is refuted in a frame iff this frame contains a subframe reducible to  $\mu$  and satisfies some other natural conditions.

M. V. Zakharyashchev [1989], *Syntax and semantics of intermediate logics*, Algebra and Logic, vol.28, pp. 402-429.

In 2005 we defined the canonical formulas for extensions of the minimal logic.

In this case we obtained a necessary and sufficient condition for the refutability of the minimal canonical formulas in general frames and proved completeness theorem for that canonical formulas.

M. Stukacheva [2006] On canonical formulas for extensions of the minimal logic // Siberian Electronic Mathematical Reports. – 2006. – Vol. 3. – P. 312–334.

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## Definition (Seegerberg, 1968)

We call Kripke  $j$ -frame, or simply  $j$ -frame, a triple  $\mu = \langle W, R, Q \rangle$ , where  $\langle W, R \rangle$  is an ordinary Kripke frame for intuitionistic logic and  $Q \subseteq W$  is a cone (upward closed set) with respect to  $R$ , which we will call the cone of abnormal worlds.

Worlds lying out of  $Q$  are called *normal*.

Note that a valuation  $v$  of a  $j$ -frame  $\mu$ , a model  $\mathcal{M}=\langle\mu, v\rangle$  and the forcing relation  $\models$  between models and formulas are defined in just the same way as for ordinary Kripke frame. The only exception is in the case of constant  $\perp$ :

$$\mathcal{M} \models_x \perp \iff x \in Q.$$

K. Segerberg [1968] *Propositional logics related to Heyting's and Johansson's*, *Theoria*, vol.34, pp.26-61.

## Definition

*A general frame for JHN is a pair  $\mathfrak{M} = \langle \mu, S \rangle$ , where  $\mu = \langle W, R, Q \rangle$  is a  $j$ -frame and  $S$  is a collection of upward closed subsets of  $W$  which contains  $\emptyset, Q$  and is closed under  $\cap, \cup$ , and operation  $\supset$  is defined as follows :*

$$\forall X, Y \in S \quad X \supset Y = \{x \in W \mid \forall y \in W (xRy \text{ and } y \in X) \Rightarrow y \in Y\}.$$

If  $S$  contains all upward closed subsets of  $W$ , then  $\mathfrak{M} = \langle \mu, S \rangle$  is in effect an ordinary Kripke  $j$ -frame.

Let  $\mathfrak{M} = \langle W, R, Q, S \rangle$  be a finite rooted general  $j$ -frame with  $e_0, \dots, e_n$  being all the distinct elements of  $W$  and  $e_0$  the root.

### Definition

A pair  $\delta = (\bar{x}, \bar{y})$  of nonempty sets  $\bar{x}, \bar{y} \subseteq W$  is called a  $d$ -domain in  $\mathfrak{M}$  if the following conditions are satisfied:

- 1  $\bar{x}$  and  $\bar{y}$  are antichains in  $W$ , and  $\bar{x}$  has at least two elements;
- 2  $(\forall x \in \bar{x})(\forall y \in \bar{y})(\neg xRy)$ ;
- 3  $(\forall z \in W)(\forall x \in \bar{x} (zRx) \Rightarrow \exists y \in \bar{y} (zRy))$ .



Let  $\mathcal{D}$  be some (possibly empty) set of  $d$ -domains in  $\mathfrak{M}$ .  
By  $\mathcal{D}_1$  we denote the set of all  $d$ -domains such that  
 $\bar{y} \cap (W \setminus Q) \neq \emptyset$  and by  $\mathcal{D}_2$  – the set of all  $d$ -domains with  
property  $\bar{y} \subseteq Q$ .

With  $\mathfrak{M} = \langle \mu, S \rangle$  and  $\mathcal{D}$  we associate the following *JHN-canonical formula*

$$J(\mu, \mathcal{D}) \Leftrightarrow (\bigwedge_{e_i R e_j} A_{ij}) \wedge (\bigwedge_{\delta \in \mathcal{D}} B_\delta) \wedge C \supset p_0,$$

where

$$C = \bigwedge_{i=1}^m (\bigwedge \Gamma_i \supset p_i \vee \perp) \supset \perp;$$

$$\Gamma_j = \{p_k \mid \neg e_j R e_k\};$$

and

- if  $\delta = (\bar{x}, \bar{y}) \in \mathcal{D}_1$ , then

$$B_\delta = \bigwedge_{e_i \in \bar{y}, e_i \notin Q} (\bigwedge \Gamma_i \supset p_i \vee \perp) \wedge \bigwedge_{e_i \in \bar{y}, e_i \in Q} (\bigwedge \Gamma_i \wedge \perp \supset p_i) \supset \bigvee_{e_j \in \bar{x}} p_j,$$

(besides it, if  $\bar{y} \cap Q = \emptyset$ , then there is no second conjunctive term);

- if  $\delta = (\bar{x}, \bar{y}) \in \mathcal{D}_2$ , then

$$B_\delta = \bigwedge_{e_i \in \bar{y}} (\bigwedge \Gamma_i \wedge \perp \supset p_i) \supset \bigvee_{e_j \in \bar{x}} p_j;$$

a term  $A_{ij}$  define as follows:

- if  $e_i \notin Q$ ,  $e_j \notin Q$ , then  $A_{ij} \equiv (\wedge \Gamma_j \supset p_j \vee \perp) \supset p_i$ ;
- if  $e_i \notin Q$ ,  $e_j \in Q$ , then  $A_{ij} \equiv (\wedge \Gamma_j \wedge \perp \supset p_j) \supset p_i$ ;
- if  $e_i \in Q$ ,  $e_j \in Q$ , then  $A_{ij} \equiv (\wedge \Gamma_j \supset p_j) \supset p_i$ .

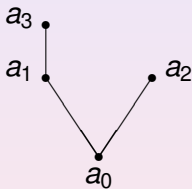
The Scott logic

$$\mathbf{SL} = \mathbf{Li} + \{(\neg\neg p \supset p) \supset p \vee \neg p\} \supset \neg p \vee \neg\neg p\}$$

(where  $\mathbf{Li}$  is the intuitionistic logic and  $\mathbf{Li} = \mathbf{Ij} + \{\perp \supset p\}$ )

is one of the first examples of an intermediate logic with the disjunction property.

Zakharyashchev proved that  $\mathbf{SL} = \mathbf{Li} + X(\mu, \mathcal{D}, \perp)$ , where  $X(\mu, \mathcal{D}, \perp)$  is the intuitionistic canonical formula with the frame  $\mu$ :



We study the paraconsistent analogue

$$\mathbf{Ls} = \mathbf{Lj} + \{(\neg\neg p \supset p) \supset p \vee \neg p\} \supset \neg p \vee \neg\neg p\}$$

of the Scott logic.

## Theorem

$$\mathbf{Ls} = \mathbf{Lj} + J(\eta_1, \mathcal{D}^1) + J(\eta_2, \mathcal{D}^2) + J(\eta_3, \mathcal{D}^3) + J(\eta_4, \mathcal{D}^4) + J(\eta_5, \mathcal{D}^5),$$

where  $J(\eta_1, \mathcal{D}^1)$ ,  $J(\eta_2, \mathcal{D}^2)$ ,  $J(\eta_3, \mathcal{D}^3)$ ,  $J(\eta_4, \mathcal{D}^4)$ ,  $J(\eta_5, \mathcal{D}^5)$  are canonical formulas [2] with



