

Provided for non-commercial research and educational use only.
Not for reproduction or distribution or commercial use.



This article was originally published in a journal published by Elsevier, and the attached copy is provided by Elsevier for the author's benefit and for the benefit of the author's institution, for non-commercial research and educational use including without limitation use in instruction at your institution, sending it to specific colleagues that you know, and providing a copy to your institution's administrator.

All other uses, reproduction and distribution, including without limitation commercial reprints, selling or licensing copies or access, or posting on open internet sites, your personal or institution's website or repository, are prohibited. For exceptions, permission may be sought for such use through Elsevier's permissions site at:

<http://www.elsevier.com/locate/permissionusematerial>

Letter to the Editor

Kaczmarz algorithm in Hilbert space and tight frames [☆]

Ryszard Szwarc

Institute of Mathematics, University of Wrocław, pl. Grunwaldzki 2/4, 50-384 Wrocław, Poland

Available online 29 November 2006

Communicated by Charles K. Chui on 10 September 2005

Abstract

We prove that any tight frame $\{g_n\}_{n=0}^{\infty}$, with $\|g_0\| = 1$, in a Hilbert space can be obtained by the Kaczmarz algorithm. The uniqueness of the correspondence is determined.

© 2006 Elsevier Inc. All rights reserved.

Keywords: Kaczmarz algorithm; Hilbert space; Bessel sequence; Tight frame

1. Introduction

Let $\{e_n\}_{n=0}^{\infty}$ be a linearly dense sequence of unit vectors in a Hilbert space \mathcal{H} . In 1937 Kaczmarz considered the problem of reconstructing vectors x from the data $\langle x, e_n \rangle$. He proved that in the finite dimensional case we have $x_n \rightarrow x$ for any x , where elements x_n are defined recursively by

$$\begin{aligned} x_0 &= \langle x, e_0 \rangle e_0, \\ x_n &= x_{n-1} + \langle x - x_{n-1}, e_n \rangle e_n. \end{aligned}$$

This formula is called the Kaczmarz algorithm [1].

It can be shown that if vectors g_n are given by the recurrence relation

$$g_0 = e_0, \quad g_n = e_n - \sum_{i=0}^{n-1} \langle e_n, e_i \rangle g_i, \quad (1)$$

then g_0 is orthogonal to g_n for any $n \geq 1$ and

$$x_n = \sum_{i=0}^n \langle x, g_i \rangle e_i. \quad (2)$$

By (1) the vectors $\{g_n\}_{n=0}^{\infty}$ are linearly dense in \mathcal{H} . Also by definition of the algorithm the vectors $x - x_n$ and e_n are orthogonal to each other. Hence

$$\|x\|^2 = \|x - x_0\|^2 + |\langle x, g_0 \rangle|^2, \quad \|x - x_{n-1}\|^2 = \|x - x_n\|^2 + |\langle x, g_n \rangle|^2, \quad n \geq 1. \quad (3)$$

[☆] Supported by European Commission Marie Curie Host Fellowship for the Transfer of Knowledge “Harmonic Analysis, Nonlinear Analysis and Probability” MTKD-CT-2004-013389 and KBN (Poland), Grant 2 P03A 028 25.

E-mail address: szwarc@math.uni.wroc.pl.

For $n \geq 1$ let S_n denote a finite dimensional operator defined by the rule

$$S_n y = \sum_{j=0}^n \langle y, e_j \rangle g_j, \quad y \in \mathcal{H}. \quad (4)$$

Observe that the formulas (1) and (2) can be restated as

$$(I - S_{n-1})e_n = g_n, \quad (5)$$

$$(I - S_n^*)x = x - x_n. \quad (6)$$

Moreover by (3) it follows that

$$\|x - x_n\|^2 = \|(I - S_n^*)x\|^2 = \|x\|^2 - \sum_{j=0}^n |\langle x, g_j \rangle|^2. \quad (7)$$

In particular

$$\sum_{n=0}^{\infty} |\langle x, g_n \rangle|^2 \leq \|x\|^2, \quad x \in \mathcal{H}. \quad (8)$$

A sequence $\{e_n\}_{n=0}^{\infty}$ is called effective if $x_n \rightarrow x$ for any $x \in \mathcal{H}$. By virtue of (7) this is equivalent to $\|x\|^2 = \sum_{n=0}^{\infty} |\langle x, g_n \rangle|^2$ for any $x \in \mathcal{H}$, which means $\{g_n\}_{n=0}^{\infty}$ is a tight frame. We refer to [2] for more information on the Kaczmarz algorithm and to [3] for the characterization of effective sequences through the Gram matrix of the sequence $\{e_n\}_{n=0}^{\infty}$.

2. Bessel sequences

Definition 1. A sequence of vectors $\{g_n\}_{n=0}^{\infty}$ in a Hilbert space \mathcal{H} will be called a Bessel sequence if (8) holds. The sequence $\{g_n\}_{n=0}^{\infty}$ will be called a special Bessel sequence if in addition $\|g_0\| = 1$.

Observe that if $\{g_n\}_{n=0}^{\infty}$ is a special Bessel sequence then substituting $x = g_0$ into (8) implies $g_n \perp g_0$ for $n \geq 1$.

Let P_n denote the orthogonal projection onto e_n^{\perp} , the orthogonal complement to the vector e_n . By [3, (1)] we have

$$I - S_n^* = P_n P_{n-1} \dots P_0, \quad (9)$$

$$I - S_n = P_0 \dots P_{n-1} P_n. \quad (10)$$

Theorem 1. For any special Bessel sequence $\{g_n\}_{n=0}^{\infty}$ in a Hilbert space \mathcal{H} there exists a sequence $\{e_n\}_{n=0}^{\infty}$ of unit vectors such that (1) holds. In other words, any special Bessel sequence can be obtained through the Kaczmarz algorithm.

Proof. We will construct the sequence $\{e_n\}_{n=0}^{\infty}$ recursively. Set $e_0 = g_0$. Assume the unit vectors e_1, \dots, e_{N-1} have been constructed such that the formula (1) holds for $n = 0, \dots, N-1$. We want to find y such that

$$(I - S_{N-1})y = g_N, \quad \|y\| = 1. \quad (11)$$

By (10) we have $(I - S_{N-1})e_{N-1} = 0$, i.e. the operator $I - S_{N-1}$ admits nontrivial kernel. Hence the solvability of (11) is equivalent to that of

$$(I - S_{N-1})y = g_N, \quad \|y\| \leq 1. \quad (12)$$

By the Fredholm alternative the equation $(I - S_{N-1})y = g_N$ is solvable if and only if g_N is orthogonal to $\ker(I - S_{N-1}^*)$. We will check that this condition holds. Let $x \in \ker(I - S_{N-1}^*)$. Then by (7) and (8) we have

$$0 = \|(I - S_{N-1}^*)x\|^2 = \|x\|^2 - \sum_{j=0}^{N-1} |\langle x, g_j \rangle|^2 \geq \sum_{j=N}^{\infty} |\langle x, g_j \rangle|^2.$$

In particular $\langle x, g_N \rangle = 0$, i.e. $g_N \perp \ker(I - S_{N-1}^*)$.

Let y denote the unique solution to

$$(I - S_{N-1})y = g_N, \quad y \perp \ker(I - S_{N-1}).$$

The proof will be complete if we show that $\|y\| \leq 1$. Again by the Fredholm alternative we have $y \in \text{Im}(I - S_{N-1}^*)$. Let $y = (I - S_{N-1}^*)x$ for some $x \in \mathcal{H}$. We may assume that $x \perp \ker(I - S_{N-1}^*)$. In particular $\langle x, g_0 \rangle = 0$, as (9) yields $g_0 \in \ker(I - S_{N-1}^*)$. By (7) we have

$$\|y\|^2 = \|(I - S_{N-1}^*)x\|^2 = \|x\|^2 - \sum_{j=1}^{N-1} |\langle x, g_j \rangle|^2.$$

On the other hand,

$$\|y\|^2 = \langle x, (I - S_{N-1})y \rangle = \langle x, g_N \rangle.$$

Therefore

$$\|y\|^2 - \|y\|^4 = \|x\|^2 - \sum_{j=1}^N |\langle x, g_j \rangle|^2 \geq 0,$$

which implies $\|y\| \leq 1$. \square

Corollary 1. For any special tight frame $\{g_n\}_{n=0}^\infty$ in a Hilbert space \mathcal{H} there exists an effective sequence $\{e_n\}_{n=0}^\infty$ of unit vectors such that (1) holds, i.e. any special tight frame can be obtained through the Kaczmarz algorithm.

For a sequence $\{e_n\}_{n=0}^\infty$ of unit vectors the special Bessel sequence $\{g_n\}_{n=0}^\infty$ is determined uniquely by (1). However a given special Bessel sequence may correspond to many sequences of unit vectors due to two reasons. First of all for certain N the dimension of the space $\ker(I - S_{N-1})$ may exceed 1. Secondly, if we fix a unit vector u in $\ker(I - S_{N-1})$ the vector e_N can be defined as $e_N = y + \lambda u$ for any complex λ number such that $|\lambda|^2 + \|y\|^2 = 1$. In what follows we will indicate properties which guarantee a one to one correspondence between $\{e_n\}_{n=0}^\infty$ and $\{g_n\}_{n=0}^\infty$.

Definition 2. A sequence of unit vectors $\{e_n\}_{n=0}^\infty$ will be called strongly redundant if the vectors $\{e_n\}_{n=N}^\infty$ are linearly dense for any N . A special Bessel sequence $\{g_n\}_{n=0}^\infty$ will be called strongly redundant if the vectors $\{g_0\} \cup \{g_n\}_{n=N}^\infty$ are linearly dense for any N .

Proposition 1. Let sequences $\{e_n\}_{n=0}^\infty$ and $\{g_n\}_{n=0}^\infty$ satisfy (1). The sequence $\{g_n\}_{n=0}^\infty$ is strongly redundant if and only if $\{e_n\}_{n=0}^\infty$ is strongly redundant and $\langle e_n, e_{n+1} \rangle \neq 0$ for any $n \geq 0$.

Proof. Assume $\{g_n\}_{n=0}^\infty$ is strongly redundant. First we will show that the kernel of $I - S_{N-1}$ is one-dimensional and thus consists of the multiples of the vector e_{N-1} (see (10)). Assume for a contradiction that $\dim \ker(I - S_{N-1}) \geq 2$. By the Fredholm alternative we get $\dim \ker(I - S_{N-1}^*) \geq 2$. Hence there exists a nonzero vector x such that $x \perp g_0$ and $(I - S_{N-1}^*)x = 0$. By (3) we obtain

$$\|x\|^2 = \sum_{n=1}^{N-1} |\langle x, g_n \rangle|^2.$$

This and the condition (8) imply that x is orthogonal to all the vectors g_0 and $\{g_n\}_{n=N}^\infty$, which contradicts the strong redundancy assumption.

Assume $\langle e_{N-1}, e_N \rangle = 0$ for some $N \geq 1$. Then by (10) we have $e_{N-1}, e_N \in \ker(I - S_N)$ which is a contradiction as the kernel is one-dimensional.

Concerning strong redundancy of $\{e_n\}_{n=0}^\infty$, assume a vector y is orthogonal to all the vectors $\{e_n\}_{n=N}^\infty$. In particular y is orthogonal to e_N . Since $\ker(I - S_N) = \mathbb{C}e_N$, by the Fredholm alternative y belongs to $\text{Im}(I - S_N^*)$. Let $y = (I - S_N^*)x$ for some $x \in \mathcal{H}$. We may assume that $x \perp g_0$ as $g_0 \in \ker(I - S_N^*)$. By (9), since y is orthogonal to e_n for

$n \geq N$, we get $y = (I - S_n^*)x = (I - S_N^*)x$ for $n \geq N$. On the other hand, by (2) and (6) we obtain that $\langle x, g_n \rangle = 0$ for $n > N + 1$. Since $x \perp g_0$, by strong redundancy assumptions we obtain $x = 0$ and thus $y = 0$.

For the converse implication assume $\{e_n\}_{n=0}^{\infty}$ is strongly redundant and $\langle e_n, e_{n+1} \rangle \neq 0$. By the inequality (see [2])

$$\|x - x_n\| \geq |\langle e_{n-1}, e_n \rangle| \|x - x_{n-1}\|$$

we get that $x - x_n \neq 0$ for any $x \perp e_0$. Since $x - x_n = (I - S_n^*)x$, the kernel of $I - S_n^*$ consists of the multiples of $e_0 = g_0$, only.

Let x be orthogonal to $\{g_0\} \cup \{g_n\}_{n \geq N+1}$ for some $N \geq 1$. By (2) we obtain that $x_n = x_N$ for $n \geq N$. By the definition of the Kaczmarz algorithm we get $x - x_N \perp e_n$ for $n \geq N + 1$. Now strong redundancy of $\{e_n\}_{n=0}^{\infty}$ implies $x - x_N = 0$. By (6) we obtain $(I - S_N^*)x = 0$. This yields $x = 0$ since the kernel is one-dimensional and consists of the multiples of g_0 . \square

For sequences $\{e_n\}_{n=0}^{\infty}$ and $\{\sigma_n e_n\}_{n=0}^{\infty}$, where σ_n are complex numbers of absolute value 1, the Kaczmarz algorithm coincides. Therefore we will restrict our attention to *admissible* sequences of unit vectors $\{e_n\}_{n=0}^{\infty}$ such that $\langle e_n, e_{n+1} \rangle \geq 0$.

Theorem 2. *Let $\{g_n\}_{n=0}^{\infty}$ be a strongly redundant special Bessel sequence. Then there exists a unique admissible sequence $\{e_n\}_{n=0}^{\infty}$ of unit vectors such that (1) holds. Moreover, the sequence $\{e_n\}_{n=0}^{\infty}$ is strongly redundant.*

Proof. The proof will go by induction. The vector e_0 is determined by $e_0 = g_0$. Assume the vectors e_0, \dots, e_{N-1} were determined uniquely. We have to show that the problem

$$(I - S_{N-1})y = g_N, \quad \|y\| = 1, \quad \langle y, e_{N-1} \rangle \geq 0$$

has a unique solution y .

By the proof of Proposition 1 the kernel of $I - S_{N-1}$ is one-dimensional and thus consists of the multiples of the vector e_{N-1} . By the proof of Theorem 1 there exists a unique solution y_N to the problem

$$(I - S_{N-1})y = g_N, \quad y \perp \ker(I - S_{N-1})$$

and $\|y_N\| \leq 1$. Moreover by this proof $\|y_N\| = 1$ if and only if

$$\|x\|^2 - \sum_{j=1}^N |\langle x, g_j \rangle|^2 = 0,$$

where $y_N = (I - S_{N-1}^*)x$ and $x \perp \ker(I - S_{N-1}^*)$. This leads to a contradiction because by inequality (8) we get that x is orthogonal to all the vectors g_0 and $\{g_n\}_{n=N}^{\infty}$. Hence $\|y_N\| < 1$.

At this stage we know that any solution to the equation

$$(I - S_{N-1})y = 0$$

is of the form

$$y = y_N + \lambda e_{N-1}, \quad \lambda \in \mathbb{C},$$

because $\ker(I - S_{N-1}) = \mathbb{C}e_{N-1}$. Since $\|y_N\| < 1$ and $y_N \perp e_{N-1}$ there exists a unique solution y satisfying $\|y\| = 1$ and $\langle y, e_{N-1} \rangle \geq 0$ namely the one corresponding to $\lambda = \sqrt{1 - \|y_N\|^2}$. \square

Corollary 2. *Let $\{g_n\}_{n=0}^{\infty}$ be a strongly redundant special tight frame. Then there exists a unique admissible effective sequence $\{e_n\}_{n=0}^{\infty}$ of unit vectors such that (1) holds. Moreover, the sequence $\{e_n\}_{n=0}^{\infty}$ is strongly redundant.*

References

- [1] S. Kaczmarz, Approximate solution of systems of linear equations, Bull. Acad. Polon. Sci. Lett. A 35 (1937) 355–357 (in German); English transl.: Internat. J. Control 57 (6) (1993) 1269–1271.
- [2] S. Kwapień, J. Mycielski, On the Kaczmarz algorithm of approximation in infinite-dimensional spaces, Studia Math. 148 (2001) 75–86.
- [3] R. Haller, R. Szwarc, Kaczmarz algorithm in Hilbert space, Studia Math. 169 (2005) 123–132.